

CS30017 编译

第六讲：类型推导

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大纲

一、类型推导问题

二、标识符作用域

三、类型约束和求解

一、类型推导问题



类型系统基本概念

- 类型系统包括由类型和规则组成
- 类型：
 - 基础类型（Primitive Type）
 - 标量类型（Scalar Types）：bool、char、int、float
 - 复合类型（Compound Type）：数组、元组
 - 自定义类型：结构体、枚举
- 类型规则：
 - 类型推导和检查规则
 - 隐式类型转换

类型系统分类



动态类型 vs 静态类型

- 静态类型系统：编译时检查类型的一致性
- 动态类型系统：运行时检查类型的一致性

```
//python代码  
def foo(x):  
    if x == 1:  
        return "bingo!"  
    return x
```

```
//foo的类型是什么?  
print(foo(10))  
print(foo(1))  
print(foo(10) + foo(1))
```

```
#: python factorial.py  
10  
bingo!  
Traceback (most recent call last):  
  File "factorial.py", line 11, in <module>  
    print(foo(10) + foo(1))  
TypeError: unsupported operand type(s) for +: 'int' and 'str'
```

强类型 vs 弱类型

- 强类型系统：一般不允许隐式类型转换
- 弱类型类型：自动隐式转换，灵活但易出错

//python代码

```
b = 1 + True;  
a = 1 + '2';  
c = '1' + True;
```

2
类型错误
类型错误

//C代码

```
int a = 1 + true;  
int b = '1' + true;  
int c = 1 + '2';  
int d = 1 + "2";
```

2
50
51
4202501

//Javascript代码

```
1 + true;  
1 + '2';  
'1' + true;
```

2
'12'
'1true'

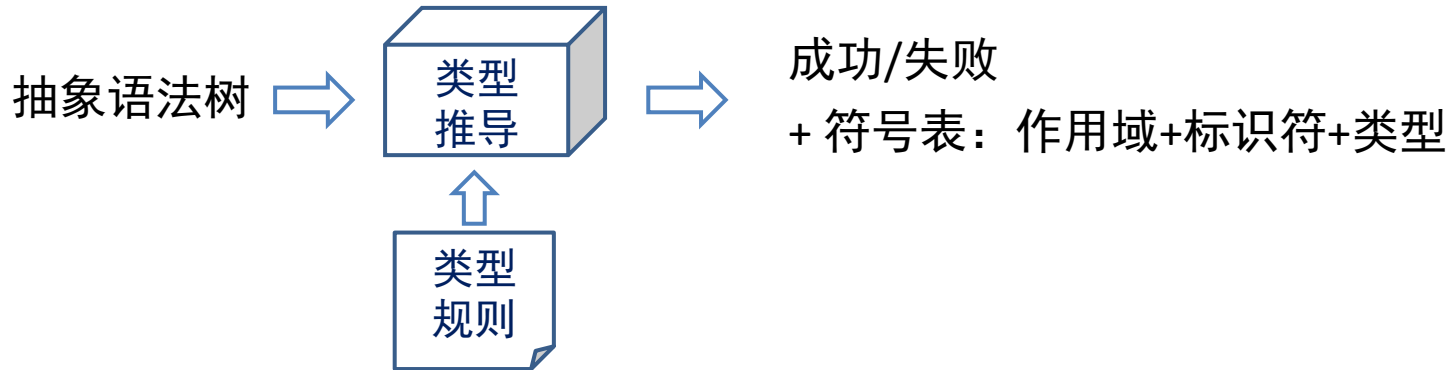
//Javascript代码

```
var a = 42;  
var b = "42";  
var c = [42];  
a === b;  
a == b;  
a == c;
```

false
true
true

类型推导问题

- 已知源代码（抽象语法树）和类型规则
- 为所有标识符找到满足类型规则的唯一解
 - 如满足运算符、函数签名等要求
- 类型检查是类型推导问题的特例



类型推导思路

- 基于抽象语法树对标识符进行作用域识别和类型分析
 - 声明新标识符：确定作用域，建立索引
 - 使用标识符：确定索引，推导或检查类型

```
let g:i32 = 10;

fn fib(x:i32) -> int {
  if (x <= 1) {
    return x;
  }
  let a = fib(x - 1);
  let b = fib(x - 2);
  let r = a + b;
  return r;
}

fn main() {
  let r = fib(10) + g;
}
```

标识符	作用域 (粗)	索引	类型
g	global	0x0000	i32
fib	global	0x0100	(i32) → i32
main	global	0x0101	(void) → void
x	fib	0x1100	i32
a	fib	0x1101	i32
b	fib	0x1102	i32
r	fib	0x1103	i32
r	main	0x82d0	i32

二、标识符作用域

作用域分析（细粒度）

```
let g:i32 = 10;

fn fib(x:i32) -> int { //scope fib: available {g, x}
    if (x <= 1) {
        return x;
    }
    let a = fib(x - 1); //{ scope 1: available {g, x, a}
        let b = fib(x - 2); //{ scope 2: available {g, x, a, b}
            let r = a + b; //{ scope 3: available {g, x, a, b, r}
                return r;
            // end scope 3 }
        // end scope 2 }
    // end scope 1 }
}

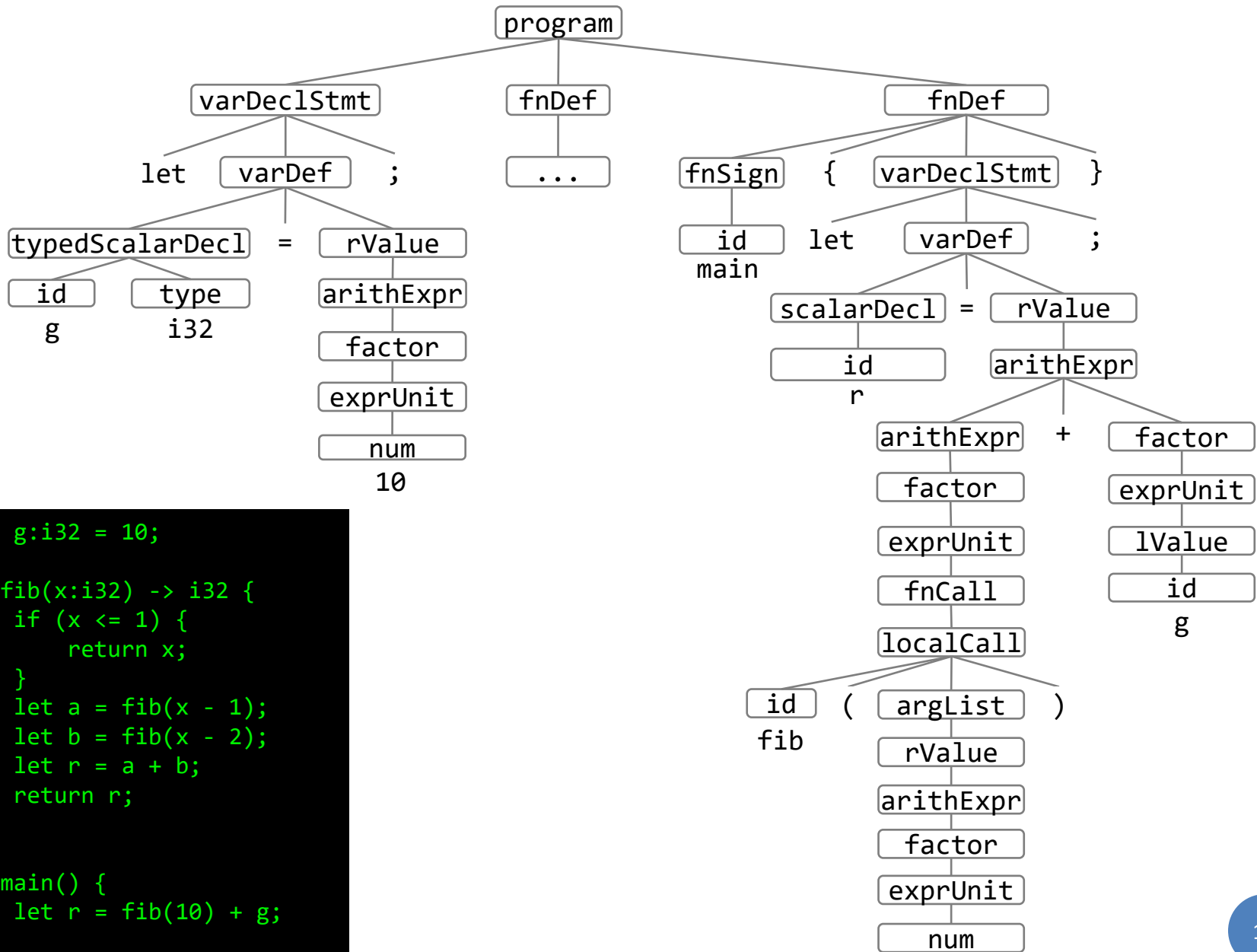
fn main() { //scope main
    let r = fib(10) + g;
}
```

作用域之间的偏序关系：fib > 1 > 2 > 3

抽象语法树： Abstract Syntax Tree

- 具体语法：程序员实际写的代码
 - 语法解析树是对源代码的完整表示
- 抽象语法树：消除解析过程中的一些步骤或节点
 - 单一展开形式塌陷，如 $E1 \rightarrow E2 \rightarrow E3 \rightarrow \text{NUM}$
 - 去除括号等冗余信息
 - 避免在叶子结点使用运算符和保留字
- AST记录编译器的阶段性分析结果，会被持续编辑

示例：Tea语言编译器生成的语法树

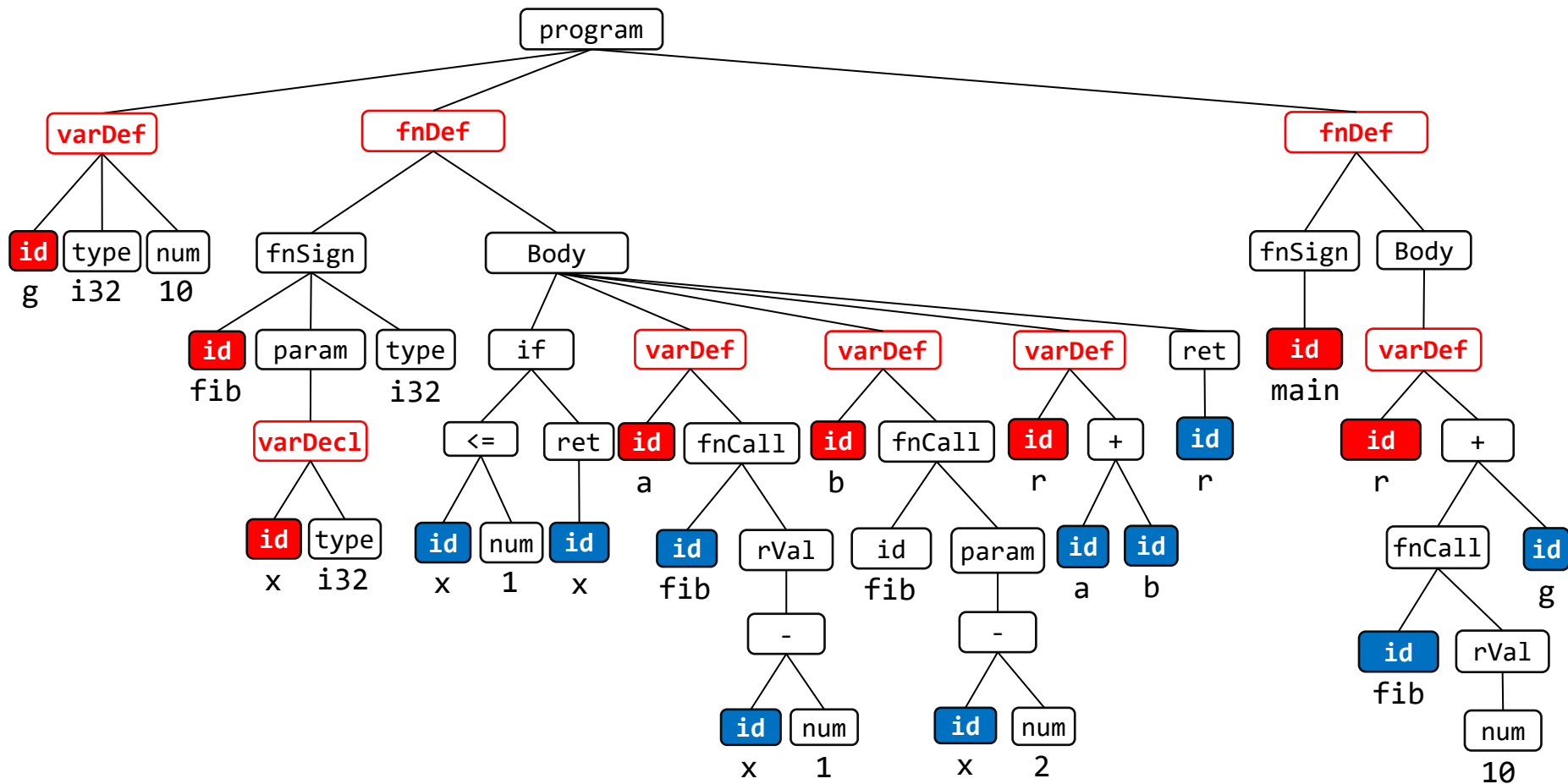


```
let g:i32 = 10;

fn fib(x:i32) -> i32 {
  if (x <= 1) {
    return x;
  }
  let a = fib(x - 1);
  let b = fib(x - 2);
  let r = a + b;
  return r;
}

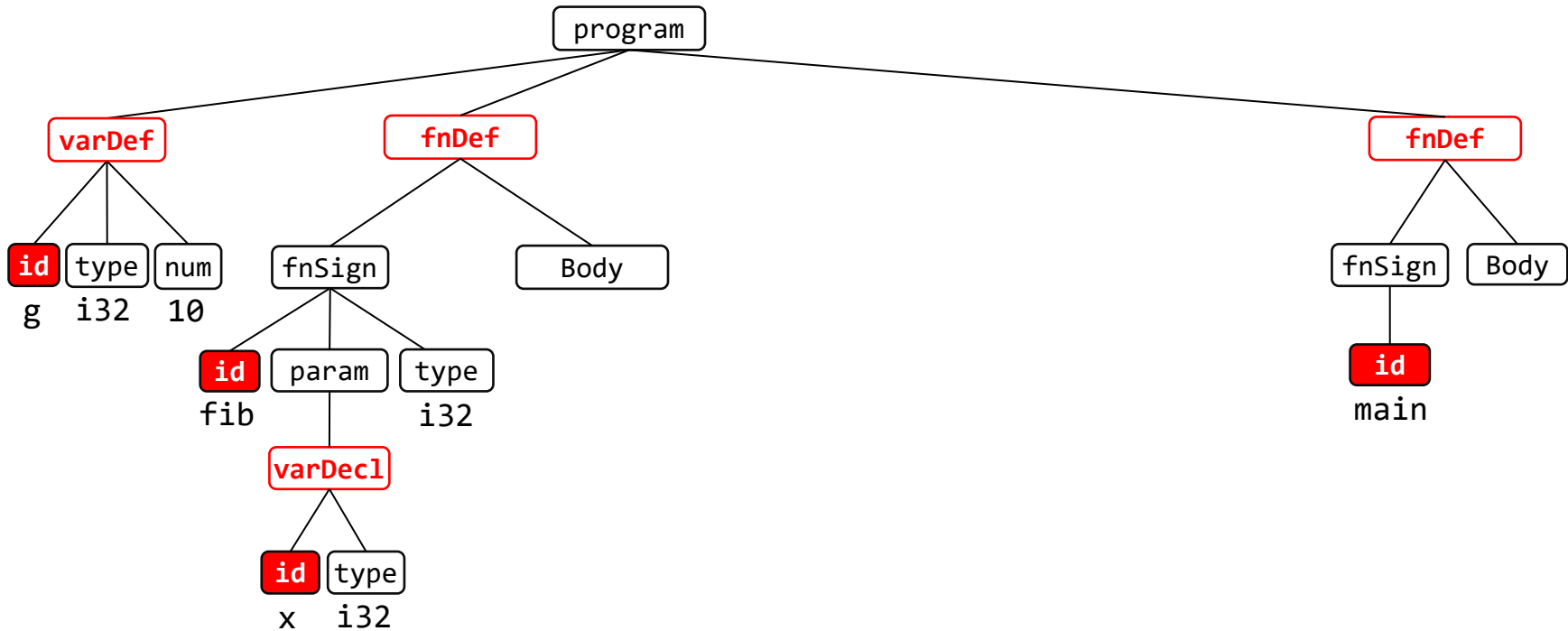
fn main() {
  let r = fib(10) + g;
}
```

语法树化简 => 问题抽象



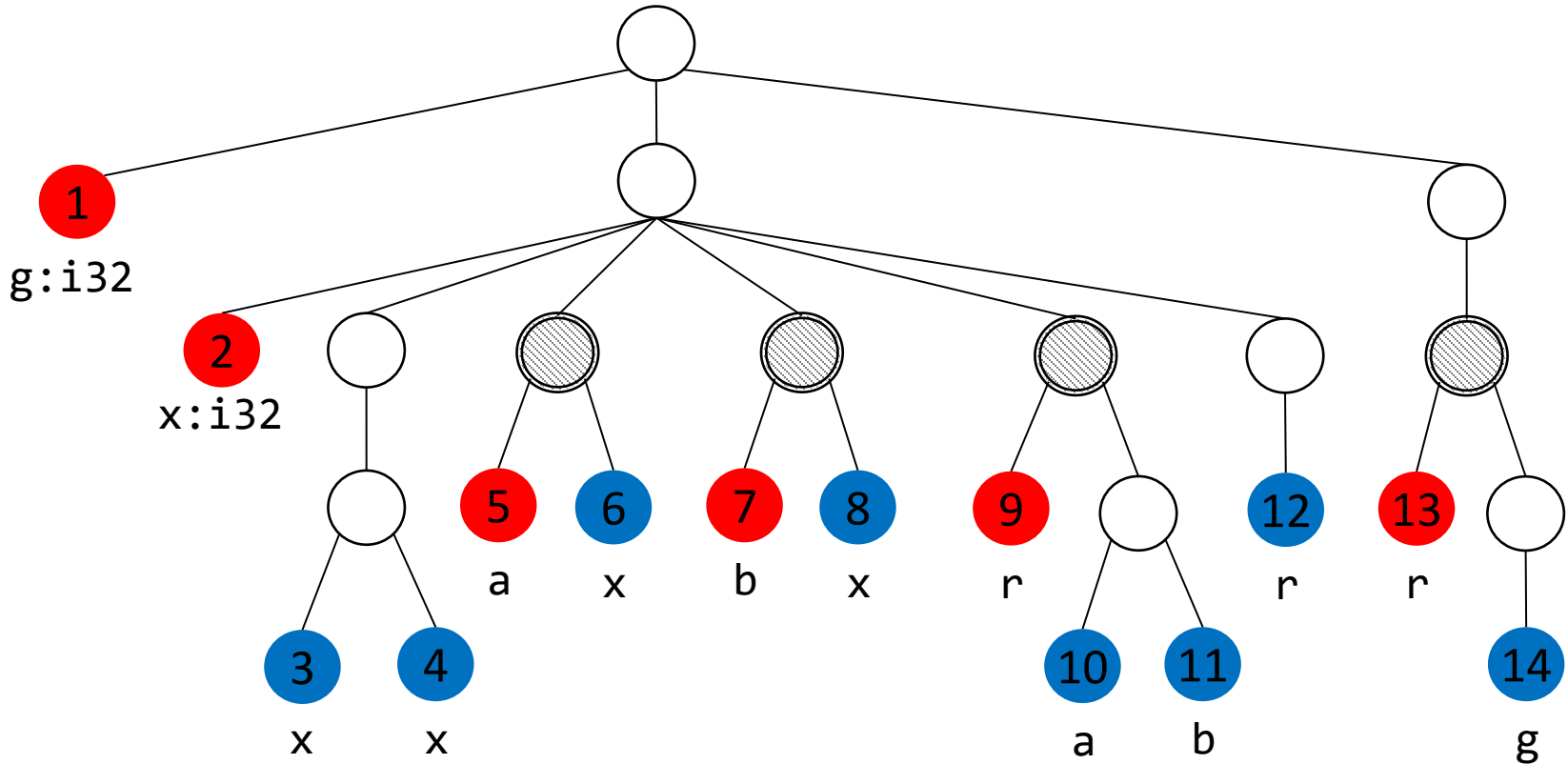
- 变量声明
- 变量引用

全局标识符符号表



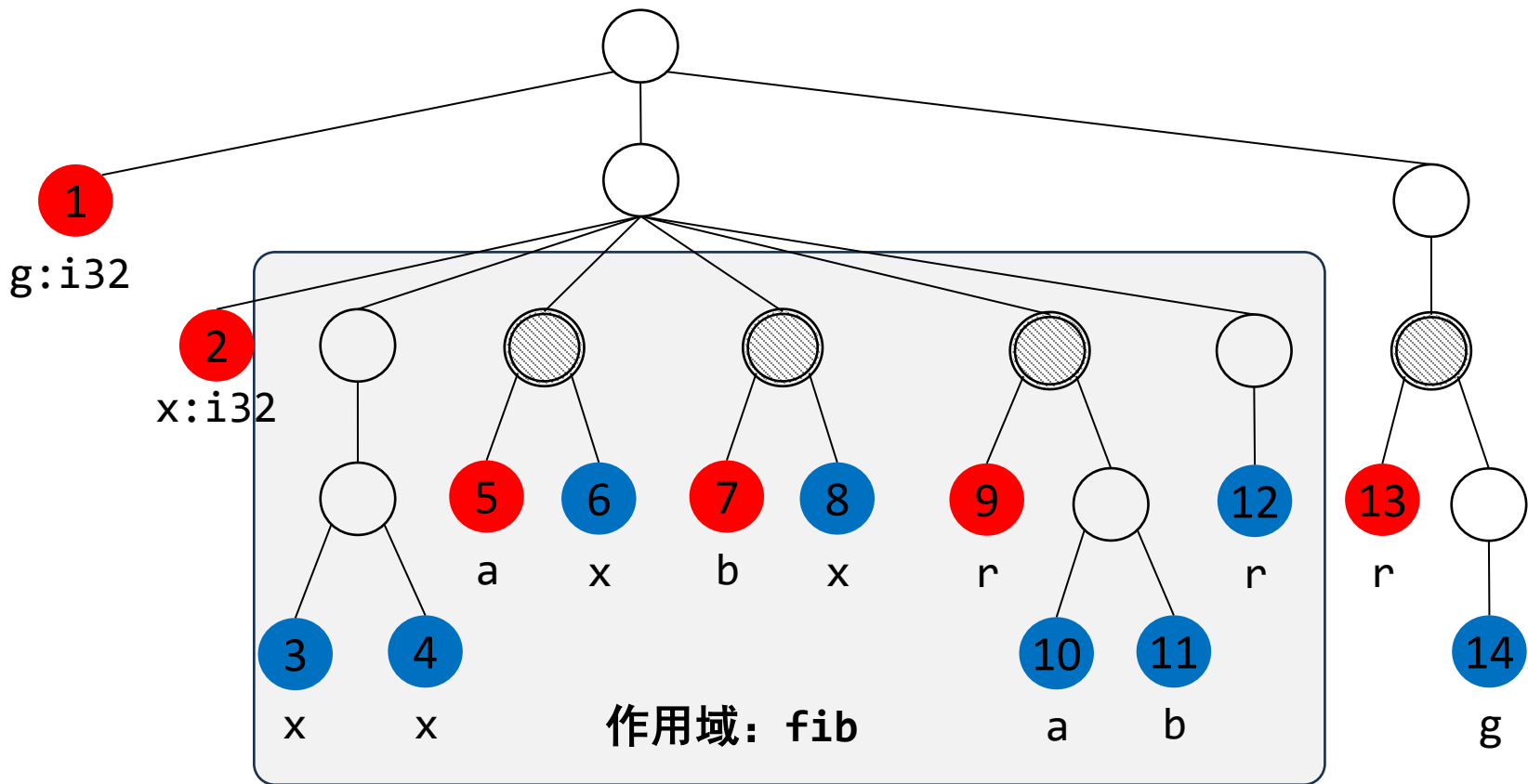
标识符	作用域	索引	类型
g	global	0x0000	i32
fib	global	0x0100	(i32) → i32
main	global	0x0101	(void) → void

局部变量类型分析



- 一般节点
- 变量声明
- 变量引用
- ⊘ 声明引用

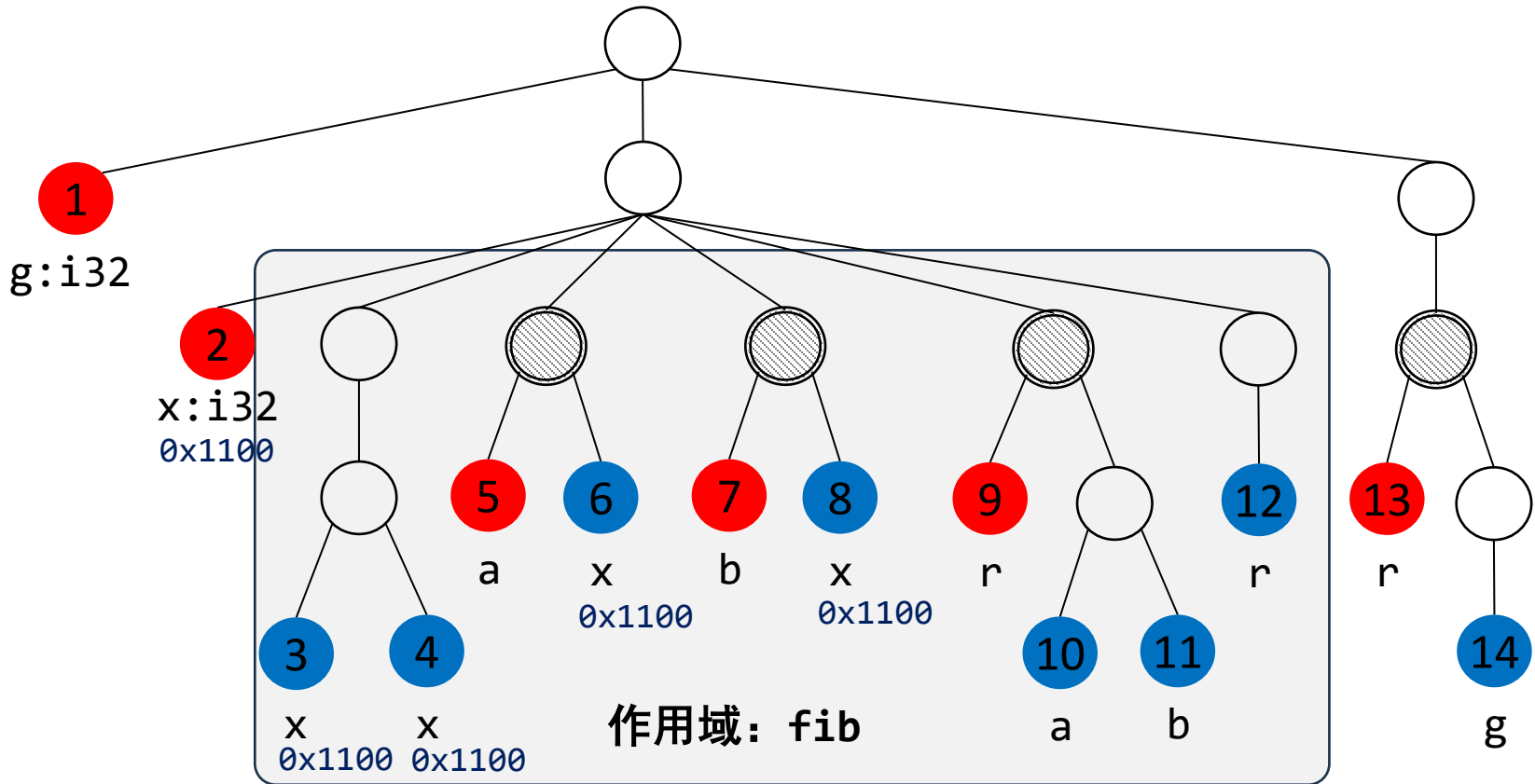
局部变量类型分析: x



标识符	作用域	索引	类型
x	fib	0x1100	i32
a			
b			
r			

- 一般节点
- 变量声明
- 变量引用
- ◐ 声明引用

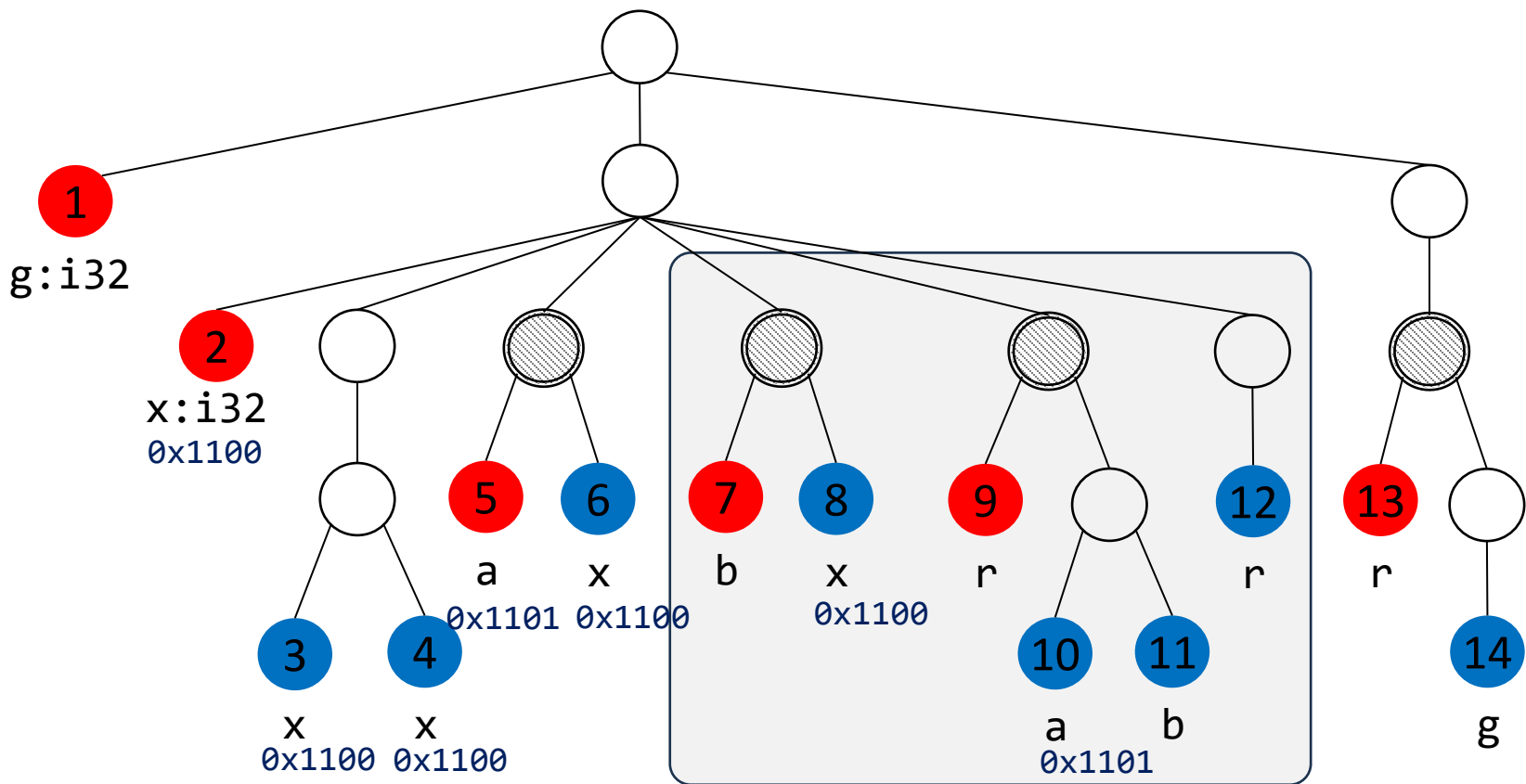
标识符索引化: x



标识符	作用域	索引	类型
x	fib	0x1100	i32
a			
b			
r			

- 一般节点
- 变量声明
- 变量引用
- ⊘ 声明引用

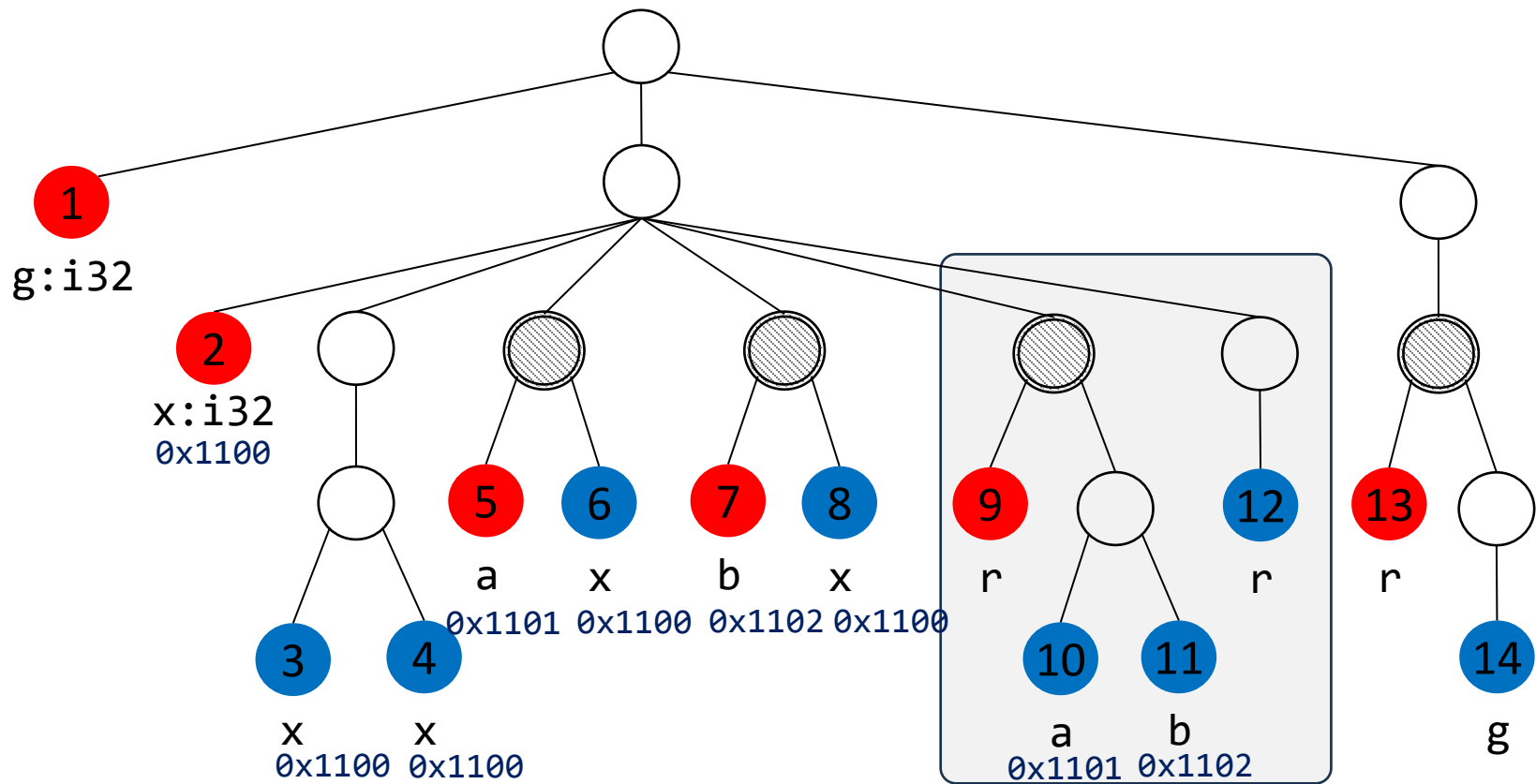
标识符索引化: a



标识符	作用域	索引	类型
x	fib	0x1100	i32
a	fib:scope1	0x1101	
b			
r			

- 一般节点
- 变量声明
- 变量引用
- ⊘ 声明引用

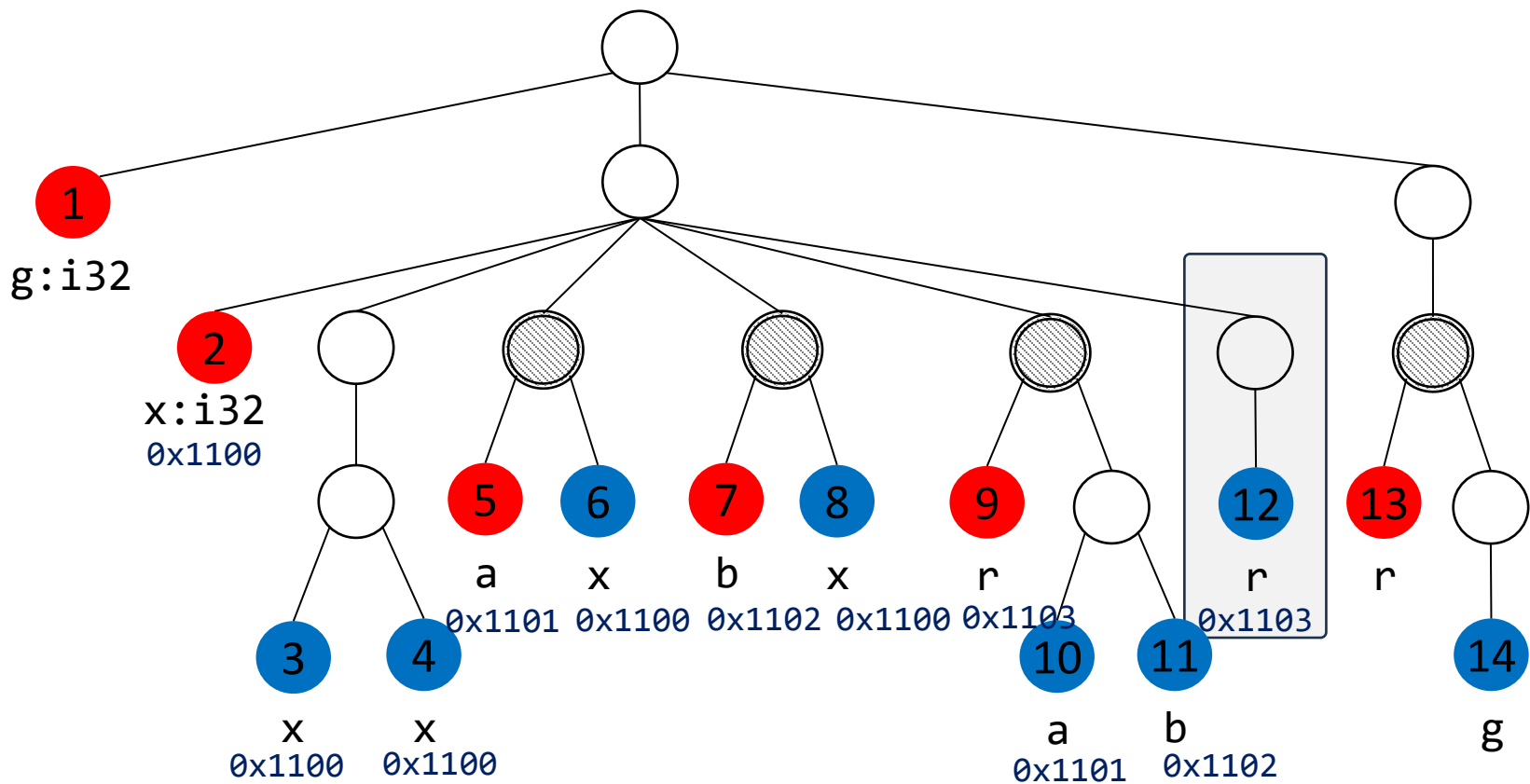
标识符索引化: b



标识符	作用域	索引	类型
x	fib	0x1100	i32
a	fib:scope1	0x1101	
b	fib:scope2	0x1102	
r			

- 一般节点
- 变量声明
- 变量引用
- ⊘ 声明引用

标识符索引化: r



标识符	作用域	索引	类型
x	fib	0x1100	i32
a	fib:scope1	0x1101	
b	fib:scope2	0x1102	
r	fib:scope3	0x1103	

- 一般节点
- 变量声明
- 变量引用
- ⊘ 声明引用

算法实现思路

- 动态维护两个哈希表记录标识符作用域和类型
 - 全局符号表
 - 局部符号表
- 遍历抽象语法树对标识符索引化
 - 标识符声明时将其加入符号表
 - 离开作用域时将其移出符号表

对应项目代码：全局变量符号表定义

```
23     /// Represents a compiled module containing all global variables and functions.
24     /// This is the top-level container for the generated IR output.
25  ✓ pub struct Module {
26     /// A map of global variable names to their definitions.
27     pub global_list: IndexMap<String, GlobalVariable>,
28     /// A map of function names to their compiled function representations.
29     pub function_list: IndexMap<String, Function>,
30 }
```

<https://github.com/tea-compiler/teac/blob/main/src/ir/module.rs>

对应项目代码：遍历全局变量

```
41  ✓      fn generate(&mut self) -> Result<(), Error> {
42          let input = self.input;
43
44          // Pass 1: handle `use` statements so imported symbols are available.
45          for use_stmt in input.use_stmts.iter() {
46              self.handle_use_stmt(use_stmt)?;
47          }
48
49          // Pass 2: register all declarations and definitions (signatures only).
50          for elem in input.elements.iter() {
51              use ast::ProgramElementInner::*;
52              match &elem.inner {
53                  VarDeclStmt(stmt) => self.handle_global_var_decl(stmt)?,
54                  FnDeclStmt(fn_decl) => self.handle_fn_decl(fn_decl)?,
55                  FnDef(fn_def) => self.handle_fn_def(fn_def)?,
56                  StructDef(struct_def) => self.handle_struct_def(struct_def)?,
57              }
58          }
59
60          // Pass 3: generate IR bodies for every function definition.
```

对应项目代码：将全局变量加入符号表

```
317  ▾      fn handle_global_var_decl(&mut self, stmt: &ast::VarDeclStmt) -> Result<(), Error> {
318          let identifier = match stmt.identifier() {
319              Some(id) => id,
320              None => return Err(Error::SymbolMissing),
321          };
322      |
323      let dtype = Dtype::try_from(stmt)?;
324      let initializers = if let ast::VarDeclStmtInner::Def(d) = &stmt.inner {
325
326      self.module
327          .global_list
328          .insert(
329              identifier.clone(),
330              GlobalVariable {
331                  dtype,
332                  identifier,
333                  initializers,
334              },
335          )
336          .map_or(Ok(()), |v| {
337              Err(Error::VariableRedefinition {
338                  symbol: v.identifier,
339              })
340          })
341      })
342  }
```

对应项目代码：函数符号表定义

```
83  ✓ pub struct FunctionGenerator<'ir> {
84      /// Shared type registry containing struct and function type definitions for
85      /// the whole module.
86      pub registry: &'ir Registry,
87      /// Reference to the module's global variable map, used during variable lookup.
88      pub global_variables: &'ir IndexMap<String, GlobalVariable>,
89      /// Map of currently visible local variables, keyed by their source-level name.
90      /// Variables are inserted when declared and removed when their enclosing scope
91      /// exits.
92      pub local_variables: IndexMap<String, LocalVariable>,
93      /// A stack of scopes. Each entry is the list of local variable names introduced
94      /// in that scope, enabling bulk removal when the scope exits via [`exit_scope`].
95      ///
96      /// [`exit_scope`]: FunctionGenerator::exit_scope
97      scope_locals: Vec<Vec<String>>,
98      /// The flat list of IR statements being accumulated for the current function.
99      pub irs: Vec<Stmt>,
100     /// The function's formal parameters as local variables.
101     pub arguments: Vec<LocalVariable>,
102     /// Counter for allocating unique virtual register indices; incremented by
103     /// [`alloc_vreg`].
104     ///
105     /// [`alloc_vreg`]: FunctionGenerator::alloc_vreg
106     pub next_vreg: usize,
107     /// Counter for allocating unique basic block label indices; starts at `1`
108     /// because index `0` is reserved for the implicit function-entry block.
109     pub next_basic_block: usize,
110 }
```

<https://github.com/tea-compiler/teac/blob/main/src/ir/function.rs>

对应项目代码：标识符作用域管理

```
176     /// Pushes a new empty lexical scope onto the scope stack.
177     ///
178     /// Call this before entering any block (e.g., `{` in the source language) so
179     /// that variables declared inside can be tracked and later removed by
180     /// [`exit_scope`].
181     ///
182     /// [`exit_scope`]: FunctionGenerator::exit_scope
183     pub fn enter_scope(&mut self) {
184         self.scope_locals.push(Vec::new());
185     }
186     |
187     /// Pops the innermost lexical scope from the scope stack and removes all local
188     /// variables that were introduced in that scope from [`local_variables`].
189     ///
190     /// [`local_variables`]: FunctionGenerator::local_variables
191     ✓ pub fn exit_scope(&mut self) {
192         if let Some(locals) = self.scope_locals.pop() {
193             for id in locals {
194                 self.local_variables.shift_remove(&id);
195             }
196         }
197     }
```

对应项目代码：while作用域处理示例

```
462 ✓      pub fn handle_while_stmt(&mut self, stmt: &ast::WhileStmt) -> Result<(), Error> {
463          let test_label = self.alloc_basic_block();
464          let true_label = self.alloc_basic_block();
465          let false_label = self.alloc_basic_block();
466
467          // Jump unconditionally into the loop test from the predecessor block.
468          self.emit_jump(test_label.clone());
469
470          // Emit the loop condition test.
471          self.emit_label(test_label.clone());
472          self.handle_bool_unit(&stmt.bool_unit, true_label.clone(), false_label.clone())?;
473
474          // Loop body; `continue` → test_label, `break` → false_label.
475          self.emit_label(true_label);
476          self.enter_scope();
477          for s in stmt.stmts.iter() {
478              self.handle_block(s, Some(test_label.clone()), Some(false_label.clone()))?;
479          }
480          self.exit_scope();
481          // Back-edge: jump back to the loop condition.
482          self.emit_jump(test_label);
483
484          self.emit_label(false_label);
485          Ok(())
486      }
```

作用域导致的类型错误举例

```
fn foo(n:i32) -> i32 {  
  while (n>0) {  
    ...  
    n = n-1;  
  }  
  return x;  
}
```

错误：变量x未声明

当前使用变量不在符号表中

```
fn foo(n:i32) -> i32 {  
  while (n>0) {  
    let x:i32;  
    ...  
    n = n-1;  
  }  
  return x;  
}
```

声明变量x，但作用域太小

错误：变量x未声明

作用域导致的类型检查问题

```
fn foo(n:i32) -> i32 {  
  let x:i32;  
  if (n>0) {  
    let x:i32;  
    ...  
    x = n-1;  
  }  
  return x;  
}
```

声明变量x

错误：重复声明

当前变量名已经在局部变量表中


```
fn foo(n:i32) -> i32 {  
  if (n>0) {  
    let x:i32;  
    ...  
  }  
  else {  
    let x:i32;  
    ...  
  }  
}
```

正确：声明变量x

正确：声明变量x


Tea语言中的全局变量使用要求

```
fn foo(n:i32) -> i32 {  
    x = x + n;  
    return x;  
}  
let x:i32 = 0;
```




全局变量声明和在函数中的使用顺序无关

```
let a:i32 = 5;  
let x:i32 = a + 5;
```



语法规则支持，翻译IR时不支持

```
let x:i32 = 0;  
fn foo(n:i32) -> i32 {  
    let x = x + n;  
    return x;  
}
```



全局变量和局部变量可以重名

对应项目代码：符号查询

```
164  ✓      pub fn lookup_variable(&self, id: &str) -> Result<Operand, Error> {
165          if let Some(local) = self.local_variables.get(id) {
166              Ok(Operand::from(local))
167          } else if let Some(global) = self.global_variables.get(id) {
168              Ok(Operand::Global(global.clone()))
169          } else {
170              Err(Error::VariableNotDefined {
171                  symbol: id.to_string(),
172              })
173          }
174      }
```

Tea语言中的函数重名问题：是否允许重载/多态？

```
fn foo(x: i32);  
fn foo() -> i32;
```



Tea语言暂时不允许重载

支持方法：在全局符号表中增添项目即可

标识符	作用域	索引	类型
foo	global	0xadc2	(i32) → void
foo	global	0xbc70	(void) → i32

对应项目代码：函数注册

```
375  ✓    fn handle_fn_decl(&mut self, decl: &ast::FnDecl) -> Result<(), Error> {
376        let identifier = decl.identifier.clone();

398        if let Some(ftype) = self
399            .registry
400            .function_types
401            .insert(identifier.clone(), function_type.clone())
402        {
403            if ftype != function_type {
404                return Err(Error::ConflictedFunction { symbol: identifier });
405            }
406        }
    }
```

三、类型约束和求解

Tea语言的类型系统

- 基础类型
 - 标量类型：i32、bool
 - 复合类型：数组
 - 函数类型
- 自定义类型：使用struct定义
- 规则（静态类型系统）
 - 类型推导：为代码中的每个标识符和表达式确定类型
 - 类型检查：分析每个参数类型是否符合运算符或函数签名要求
 - 类型转换：允许隐式类型转换？

类型问题举例

```
fn foo(n:i32) -> i32;
```

```
fn main() {  
    foo(foo);  
}
```

参数类型错误



```
fn main() {  
    let a;  
    let b;  
    b = a;  
}
```

无法确定a和b的类型



类型推导/检查思路

Damas–Hindley–Milner方法

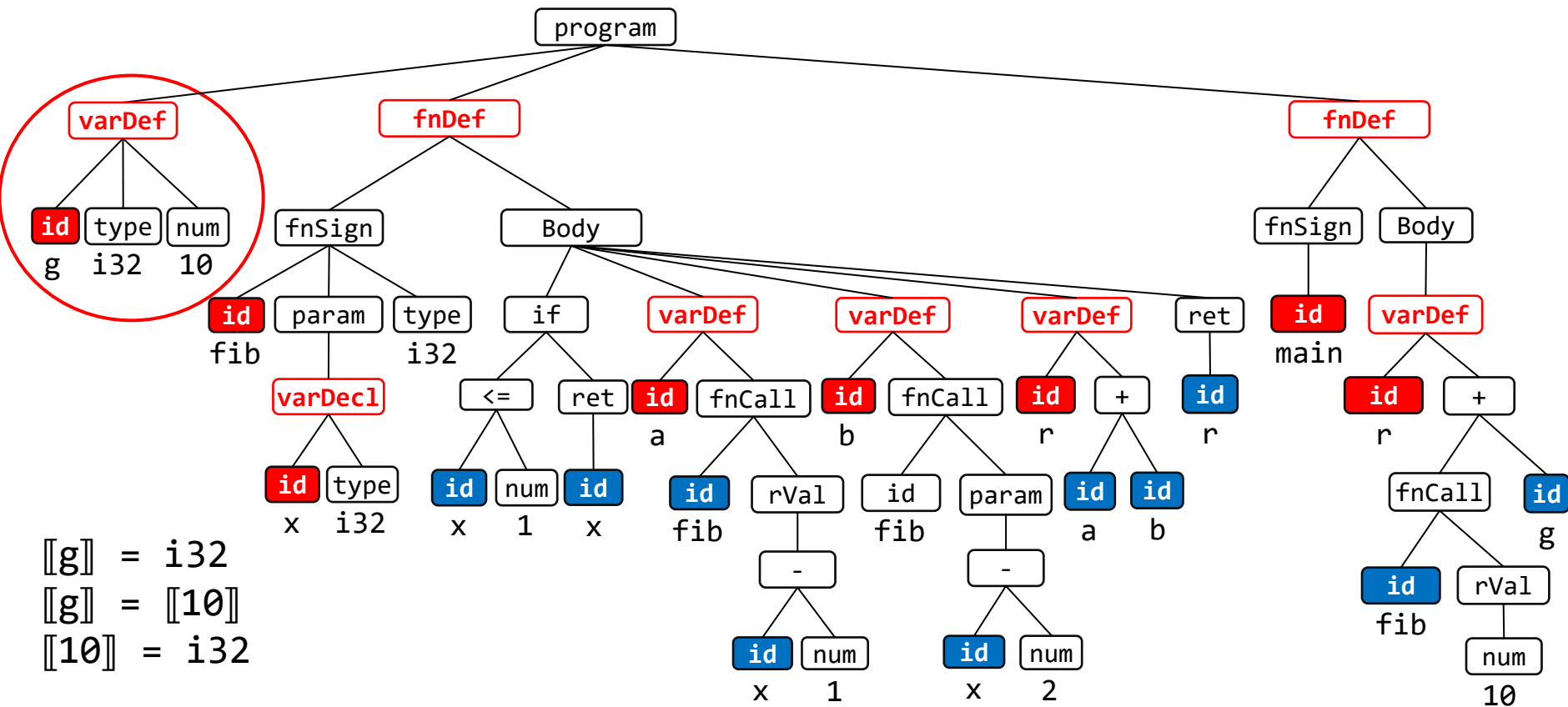
- 根据符号表确定变量类型
- 基于类型规则提取类型约束
 - 类型表示：用[[X]]表示变量X的类型
 - 约束提取：一般都为等价关系，如果支持子类型和范型除外
- 约束求解

类型规则设计

- 为不同的语法制定相应的推断规则

代码示例	代码模式	约束
<code>a: i32</code>	$X: Ty$	$\llbracket X \rrbracket = Ty$
<code>0</code>	N	$\llbracket N \rrbracket = i32$
<code>a = b;</code>	$X = Y$	$\llbracket X \rrbracket = \llbracket Y \rrbracket$
<code>a + b;</code>	$X \text{ bop } Y$	$\llbracket X \rrbracket = \llbracket Y \rrbracket = \llbracket X+Y \rrbracket$
<code>c = a + b;</code>	$Z = X \text{ bop } Y$	$\llbracket X \rrbracket = \llbracket Y \rrbracket = \llbracket X+Y \rrbracket = \llbracket Z \rrbracket$

示例：约束提取

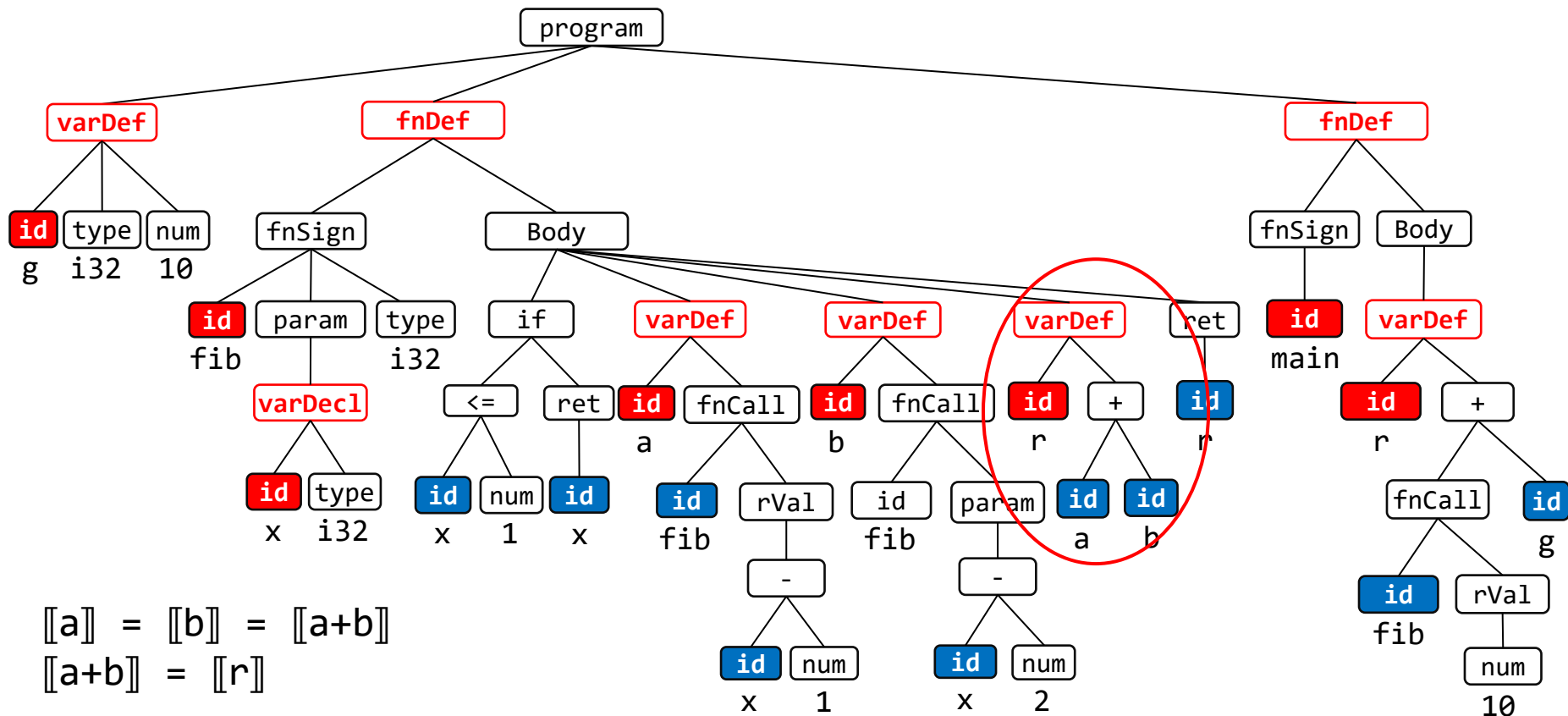


[[g]] = i32
 [[g]] = [[10]]
 [[10]] = i32

标识符	作用域	索引	类型
g	global	0x0000	i32
fib	global	0x0100	(i32) → i32
main	global	0x0101	(void) → void

● 变量声明
 ● 变量引用

示例：约束提取



标识符	作用域	索引	类型
x	fib	0x1100	i32
a	fib:scope1	0x1101	
b	fib:scope2	0x1102	
r	fib:scope3	0x1103	

- 变量声明
- 变量引用

更多类型规则

代码示例

```
a > b
```

```
a && b
```

```
if(a){...}
```

```
while(a){...}
```

代码模式

```
X cmp Y
```

```
X lop Y
```

```
if(X)
```

```
while(X)
```

约束

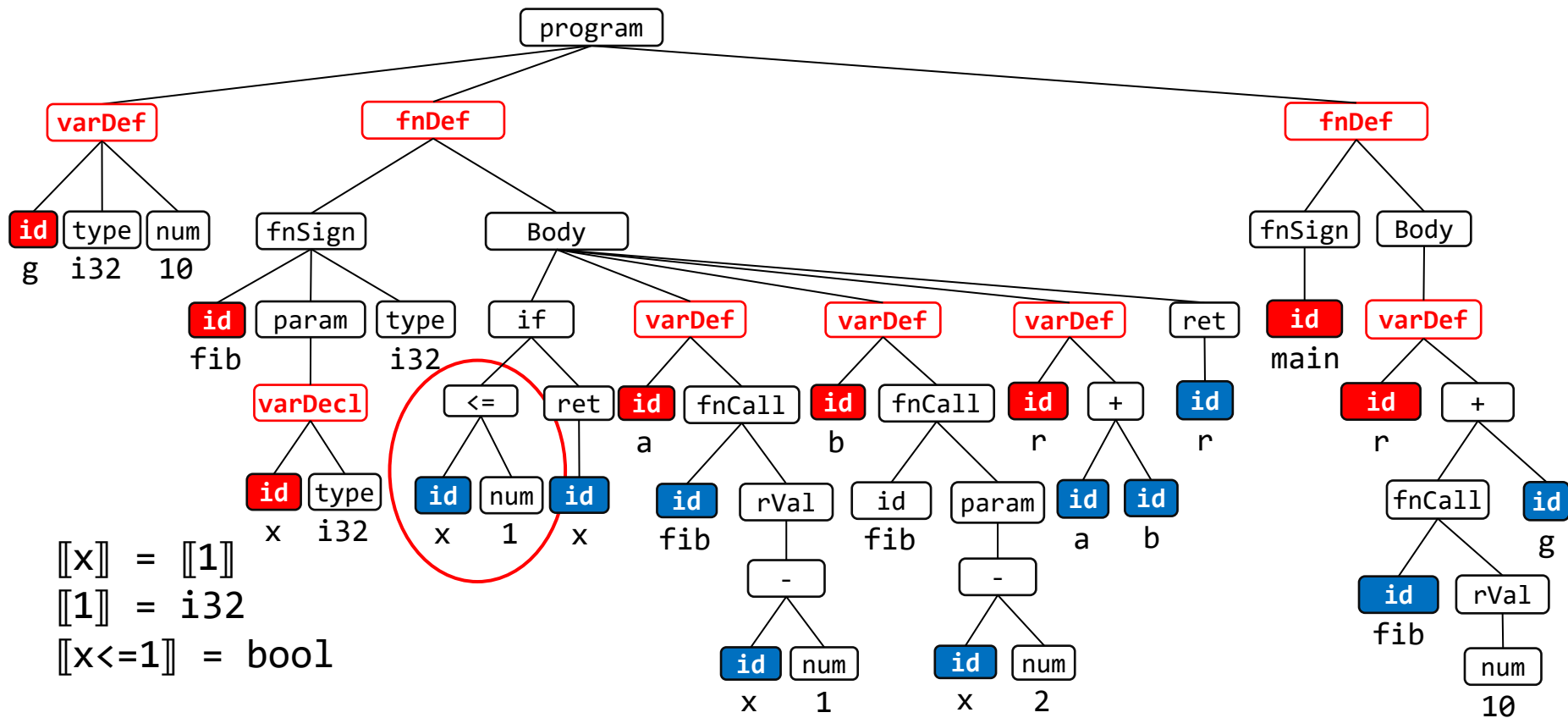
$\llbracket X \rrbracket = \llbracket Y \rrbracket$, $\llbracket X \text{ cmp } Y \rrbracket = \text{bool}$

$\llbracket X \rrbracket = \llbracket Y \rrbracket = \llbracket X \text{ lop } Y \rrbracket = \text{bool}$

$\llbracket X \rrbracket = \text{bool}$

$\llbracket X \rrbracket = \text{bool}$

示例：约束提取



标识符	作用域	索引	类型
x	fib	0x1100	i32
a	fib:scope1	0x1101	
b	fib:scope2	0x1102	
r	fib:scope3	0x1103	

- 变量声明
- 变量引用

更多类型规则

代码示例

```
foo(a, b);
```

代码模式

```
F(X, Y)
```

约束

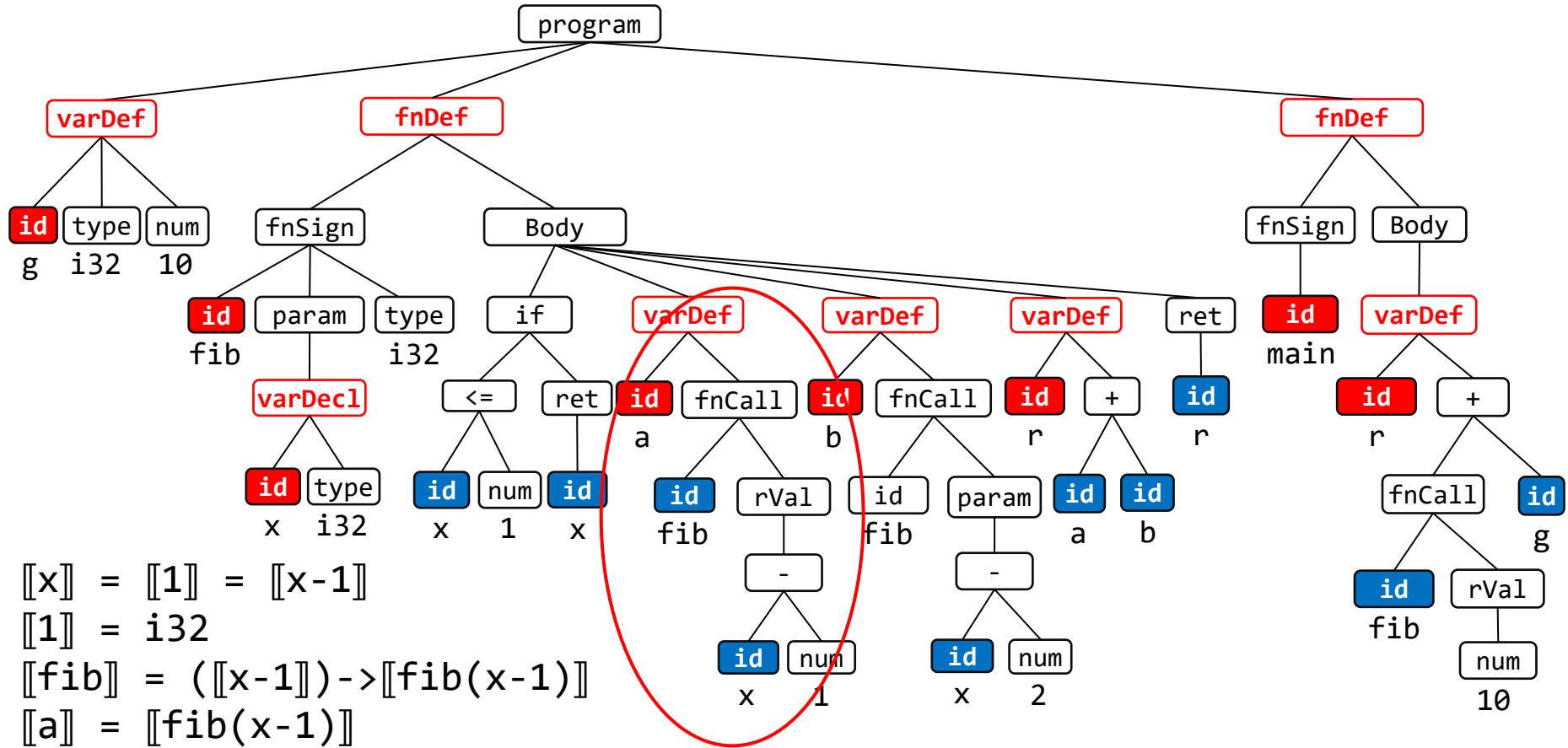
```
[[F]] = ([[X]], [[Y]]) -> [[F(X, Y)]]
```

```
return a;
```

```
F(X) -> Ty {  
    ...  
    return Y;  
}
```

```
[[F]] = ([[X]])->Ty, [[Y]] = Ty
```

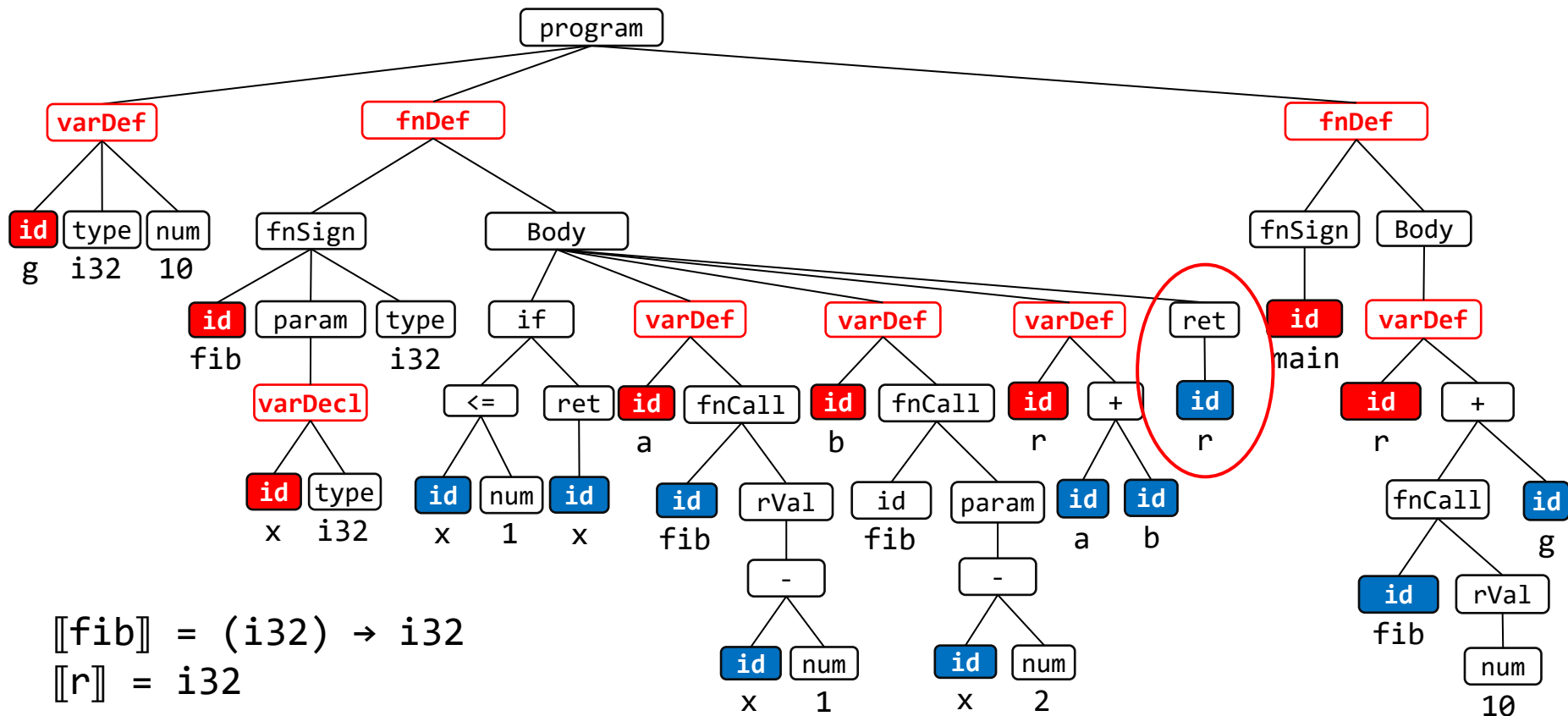
示例：约束提取



$[[x]] = [[1]] = [[x-1]]$
 $[[1]] = i32$
 $[[fib]] = ([[x-1]] \rightarrow [[fib(x-1)])]$
 $[[a]] = [[fib(x-1)]]$

标识符	作用域	索引	类型
g	global	0x0000	i32
fib	global	0x0100	(i32) → i32
main	global	0x0101	(void) → void

示例：约束提取



标识符	作用域	索引	类型
g	global	0x0000	i32
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main	global	0x0101	(void) → void

- 变量声明
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约束提取结果

全局变量:

```
[[g]] = i32
[[g]] = [[10]]
[[10]] = i32
```

main函数:

```
[[main]] = (void)->void
[[fib]] = ([[10]])->[[fib(10)]]
[[fib(10)]] = [[g]] = [[r_main]]
```

```
let g:i32 = 10;

fn fib(x:i32) -> i32 {
  if (x <= 1) {
    return x;
  }
  let a = fib(x - 1);
  let b = fib(x - 2);
  let r = a + b;
  return r;
}

fn main() {
  let r = fib(10) + g;
}
```

fib函数:

```
[[fib]] = ([[x]])->i32
[[r_fib]] = i32
[[x]] = i32

[[x]] = [[1]]
[[1]] = i32
[[x<=1]] = bool
[[x]] = i32

[[x]] = [[1]] = [[x-1]]
[[1]] = i32
[[fib]] = ([[x-1]])->[[fib(x-1)]]
[[a]] = [[fib(x-1)]]

[[x]] = [[2]] = [[x-2]]
[[2]] = i32
[[fib]] = ([[x-2]])->[[fib(x-2)]]
[[b]] = [[fib(x-2)]]

[[a]] = [[b]] = [[a+b]] = [[r_fib]]
```

基于并查集方法求解

- 维护不存在相交关系的集合，支持查找和联合两种操
 - Find(x): 返回包含变量x的集合
 - Union(x, y): 联合包含x和y的两个集合

```
while(getPair()!=NULL){  
    [p,q] = readPair(p,q);  
    pset = find(p);  
    qset = find(q);  
    if(pset == qset)  
        continue;  
    else union(p,q);  
}
```

应用并查集方法求解

$[[g]] = i32$
 $[[g]] = [[10]]$
 $[[10]] = i32$

$S1: \{i32, [[10]], [[g]]\}$

$[[fib]] = ([[x]]) \rightarrow i32$
 $[[r_{fib}]] = i32$
 $[[x]] = i32$

$S1: \{i32, [[10]], [[g]], [[x]], [[r_{fib}]]\}$

$S2: \{[[fib]], ([[x]]) \rightarrow i32\}$

$[[x]] = [[1]]$
 $[[1]] = i32$
 $[[x \leq 1]] = \text{bool}$
 $[[x]] = i32$

$S1: \{i32, [[10]], [[1]], [[g]], [[x]], [[r_{fib}]]\}$

$S3: \{\text{bool}, [[x \leq 1]]\}$

$[[x]] = [[1]] = [[x-1]]$
 $[[1]] = i32$
 $[[fib]] = ([[x-1]]) \rightarrow [[fib(x-1)]]$
 $[[a]] = [[fib(x-1)]]$

$S1: \{i32, [[10]], [[1]], [[g]], [[x]], [[r_{fib}]], [[x-1]]\}$

$S2: \{[[fib]], ([[x]]) \rightarrow i32, ([[x-1]]) \rightarrow [[fib(x-1)]]\}$

$S4: \{[[a]], [[fib(x-1)]]\}$

...

...

应用并查集方法求解

$[[x]] = [[2]] = [[x-2]]$

$[[2]] = i32$

$[[fib]] = ([[x-2]]) \rightarrow [[fib(x-2)]]$

$[[b]] = [[fib(x-2)]]$

S1:{i32, [[10]], [[1]], [[2]], [[g]], [[x]], [[r_{fib}]],
[[x-1]], [[x-2]]}

S2:{[[fib]], ([[x]]) \rightarrow i32, ([[x-1]]) \rightarrow [[fib(x-1)]],
([[x-2]]) \rightarrow [[fib(x-2)]] }

S4:{[[a]], [[b]], [[fib(x-1)]], [[fib(x-2)]]}

$[[a]] = [[b]] = [[a+b]] = [[r_{fib}]]$

S1=S1US4:

{i32, [[10]], [[1]], [[2]], [[g]], [[x]], [[r_{fib}]], [[x-1]],
[[x-2]], [[a]], [[b]], [[fib(x-1)]], [[fib(x-2)]]}

$[[main]] = (void) \rightarrow void$

S5:{(void) \rightarrow void, [[main]]}

$[[fib]] = ([[10]]) \rightarrow [[fib(10)]]$

$[[fib(10)]] = [[g]] = [[r_{main}]]$

S1:{i32, [[10]], [[1]], [[2]], [[g]], [[x]], [[r_{fib}]],
[[r_{main}]], [[x-1]], [[x-2]], [[a]], [[b]], [[fib(10)]],
[[fib(x-1)]], [[fib(x-2)]]}

S2:{[[fib]], ([[x]]) \rightarrow i32, ([[x-1]]) \rightarrow [[fib(x-1)]],
([[x-2]]) \rightarrow [[fib(x-2)]], ([[10]]) \rightarrow [[fib(10)]] }

最终解

S1:{i32, [[10]], [[1]], [[2]], [[g]], [[x]], [[r_{fib}]], [[r_{main}]], [[x-1]], [[x-2]], [[a]], [[b]], [[fib(10)]], [[fib(x-1)]], [[fib(x-2)]]}

S2:{[[fib]], ([[x]])->i32, ([[x-1]])->[[fib(x-1)]], ([[x-2]])->[[fib(x-2)]], ([[10]])->[[fib(10)]] }

S3:{bool, [[x<=1]]}

S5:{(void)->void, [[main]]}



[[g]] = i32
[[a]] = i32
[[b]] = i32
[[r_{fib}]] = i32
[[r_{main}]] = i32
[[x<=1]] = bool
[[fib]] = (i32) → i32
[[main]] = (void)->void
[[fib(T3)]] = i32

更多类型规则：数组

代码示例

代码模式

约束

```
{0, 1}
```

```
{M, N}
```

```
[[{M, N}]] = &i32
```

```
{0; 1}
```

```
{M; N}
```

```
[[{M; N}]] = &i32
```

```
let a[10]: i32
```

```
X[I]: Ty
```

```
[[X]] = &Ty
```

```
b = a[i];
```

```
Y = X[Z]
```

```
[[Z]] = i32, [[Y]] = [[*X]], [[X]] = &[[*X]]
```

```
a[i] = b;
```

```
X[Z] = Y
```

```
[[Z]] = i32, [[X]] = &[[Y]]
```

更多类型规则：结构体

代码示例

```
struct Foo {  
    a: i32,  
    b: i32,  
}
```

```
foo.a = d;
```

代码模式

```
struct ST {  
    A: Ty1,  
    B: Ty2,  
}
```

```
X.A = Y
```

约束

$$\llbracket \text{ST} \rrbracket = \llbracket \text{A}, \text{B} \rrbracket = (\text{Ty1}, \text{Ty2})$$
$$\llbracket \text{X.A} \rrbracket = \llbracket \text{Y} \rrbracket, \llbracket \text{X.A}, _ \rrbracket = \llbracket \text{X} \rrbracket$$

类型推断可能存在的问题


- 解不唯一的情况：优先选择哪些类型？
- 无解的情况：是否允许隐式类型转换？
- 如何判断类型是否等价：名字相同 vs 结构相同

```
struct Position {  
    x:i32,  
    y:i32  
}
```

```
struct Location {  
    x:i32,  
    y:i32  
}
```


递归问题

```
// Tea语言代码
fn fac(n:i32) -> i32 {
  if (n == 0) {
    return 1;
  } else {
    let r = n * fac(n-1);
    return r;
  }
}
```



[[fac]] = (i32)->i32

```
// Tea语言代码
struct List {
  v:i32,
  next List,
}
```



[[List]] = ϕ = (i32, ϕ)

```
// C代码
struct List {
  i32 data;
  struct List* next;
};
```



[[List]] = ϕ = (i32, & ϕ)
= (i32, usize)

缺少返回语句

```
fn foo(n:i32) -> i32 {  
    if (n <= 1) {  
        return n;  
    }  
    let x:i32 = foo(n-1);  
}
```

类型为i32/void时自动添加

思考

- 假如Tea语言函数声明可缺省类型，设计方法分析下列代码中的类型信息

```
fn f(n) {  
    if (n == 0) {  
        return 1;  
    } else {  
        let r = n * f(n-1);  
        return r;  
    }  
}
```

```
fn f(f, n) {  
    if (n == 0) {  
        return 1;  
    } else {  
        let r = n * f(f, n-1);  
        return r;  
    }  
}
```